University of South Australia Formula Sheet for Source and Channel Coding

Source Entropy: $H(X) = \sum_{i=1}^{K} P(X=i) \log_2(1/P(X=i))$

Mutual Information: $I(X;Y) = \sum_{i} \sum_{j} P(X=i;Y=j) \log_2 \frac{P(X=i/Y=j)}{P(X=i)}$

Conditional Entropy: $H(X/Y) = \sum_{i} \sum_{j} P(X=i;Y=j) \log_2 \frac{1}{P(X=i/Y=j)}$

Kraft Inequality: $\sum_{i=1}^K D^{-n_i} \leq 1$ Average Message Length: $L = \sum_{i=1}^K p_i n_i$

PCM μ -Law Characteristic: $\frac{c(x)}{x_{max}} = \frac{\ln(1+\mu|x|/x_{max})}{\ln(1+\mu)} \operatorname{sign}(x)$

Baseband Error Probability: Baseband $p_e = Q(A/\sigma)$, Bandpass $p_e = Q(\sqrt{(2E_b/N_0)})$ Boltzman's constant: $k = 1.3710^{-23}$ J/deg

Capacity: continuous channel $C = B \log_2(1 + S/N)$, discrete system: $\frac{E_b}{N_0} > \frac{B}{r_b}(2^{\frac{r_b}{B}} - 1)$

Block Code Generator: $c_i = d_i G$ where $G = [I_k | P]$, where d_i is the i-th row vector of information bits etc; Syndrome Decoding: $s_i^T = H r_i^T$ where $H = [P^T | I_{n-k}]$

Binomial Distribution: Let the discrete random variable I be the number of times an event E occurs in n independent trials. If P(E) = q then

$$P_I(i) = \binom{n}{i} q^i (1-q)^{n-i}$$

where $\binom{n}{i} = \frac{n!}{i!(n-i)!}$

